





FFMK: A FAST AND FAULT-TOLERANT MICROKERNEL-BASED SYSTEM FOR EXASCALE COMPUTING

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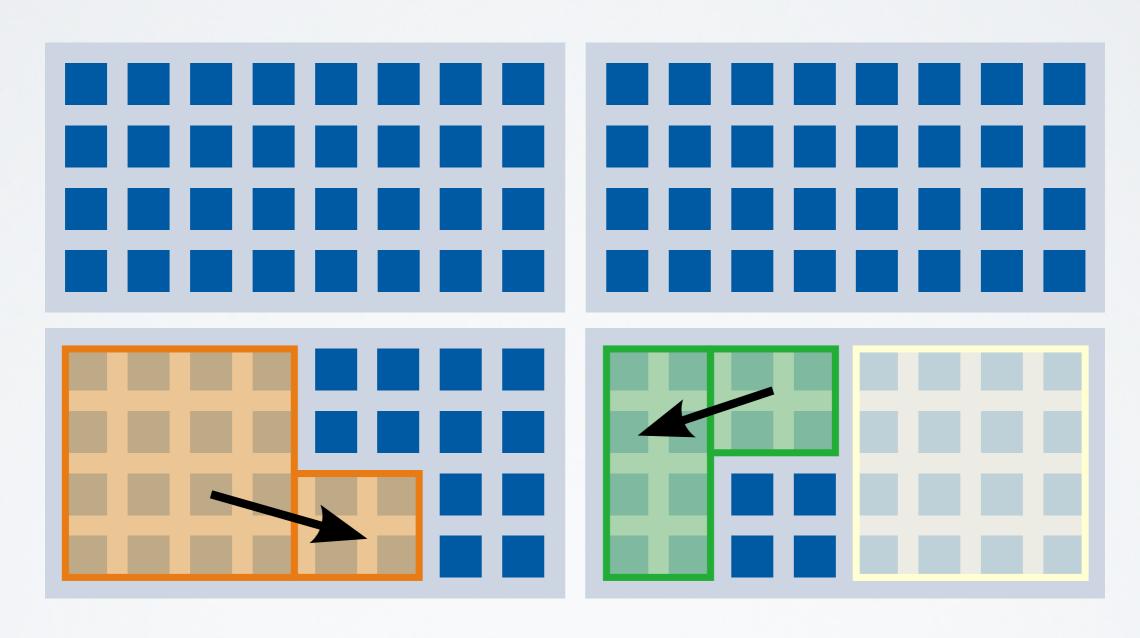
CARSTEN WEINHOLD, TU DRESDEN







SYSTEM MODEL



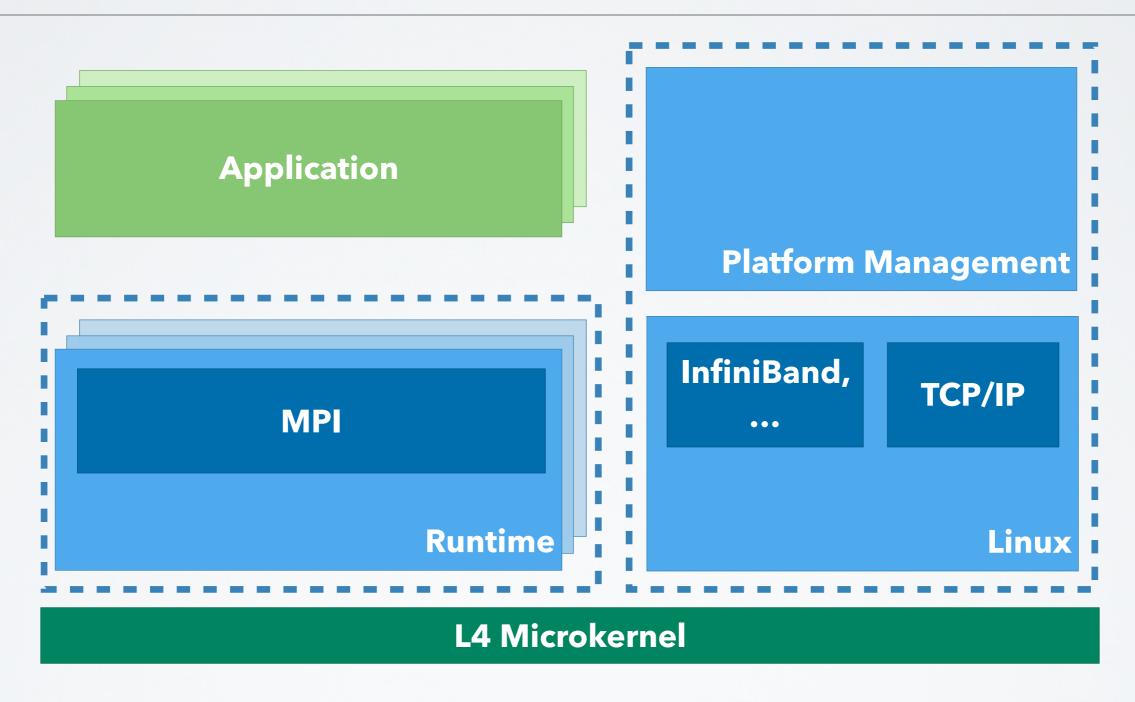








NODE ARCHITECTURE





- L4 microkernel controls the node
- Light-weight and low-noise
- Virtualization: L⁴Linux on L4 microkernel
- Unmodified Linux programs (MPI, ...)
- Linux process = L4 task + L4 threads
- Linux syscalls / exceptions:
 generic forwarding to L⁴Linux kernel



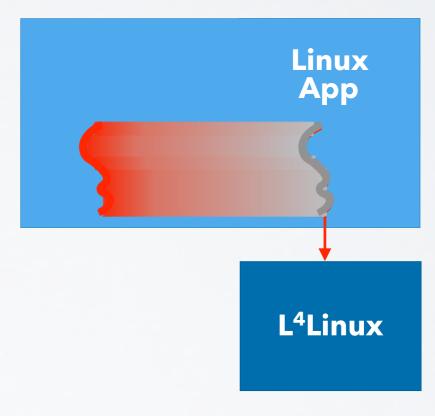
L4 Microkernel / Hypervisor

Core Core Core Core Core



DECOUPLED EXECUTION

Decoupling: move Linux thread to new L4 thread on its own core

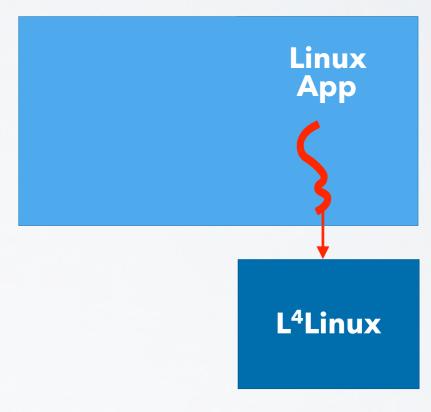






DECOUPLED EXECUTION

- Decoupling: move Linux thread to new L4 thread on its own core
- Linux syscall: Move back to Linux

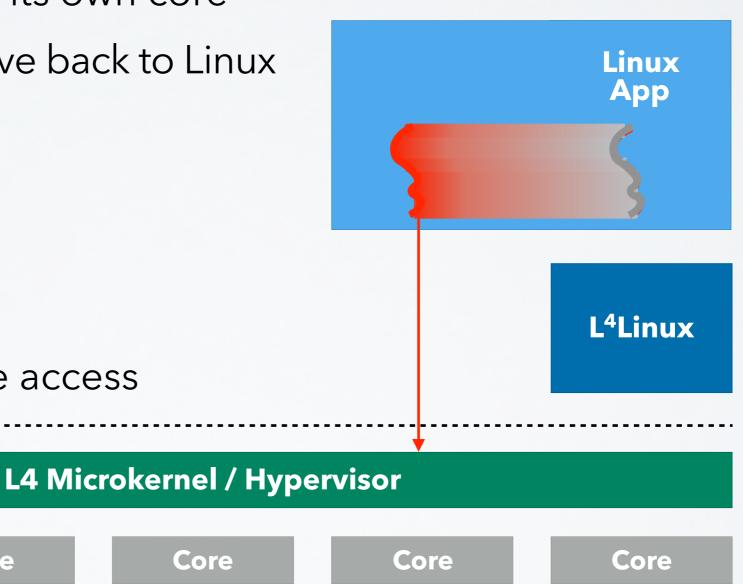


Core Core Core Core Core



DECOUPLED EXECUTION

- **Decoupling:** move Linux thread to new L4 thread on its own core
- Linux syscall: Move back to Linux
- L4 syscalls:
 - Scheduling
 - Threads
 - Memory
- Direct I/O device access



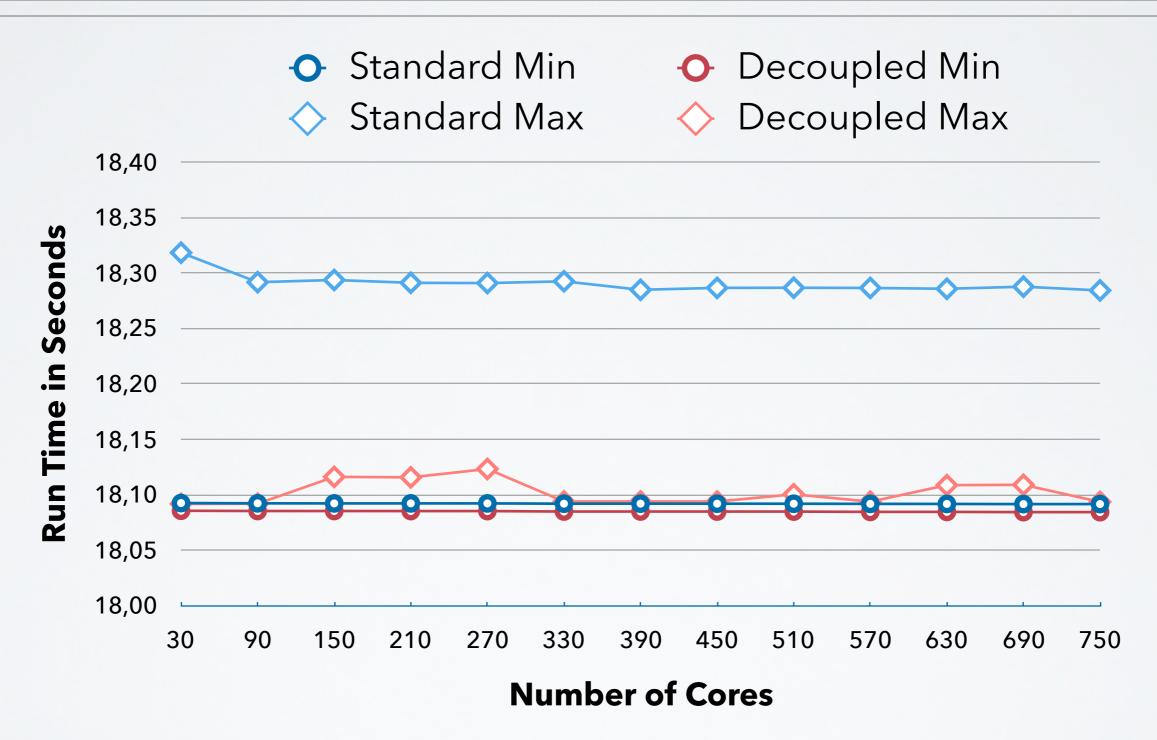
Core Core Core







DECOUPLING: EP



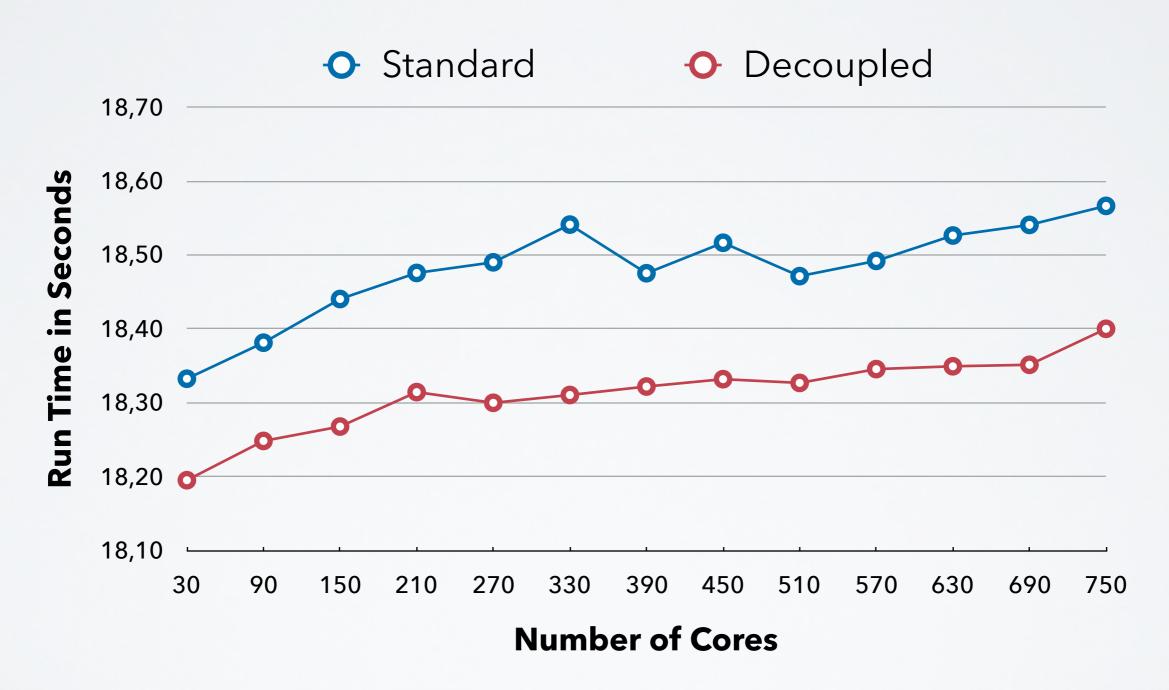
Adam Lackorzynski, Carsten Weinhold, Hermann Härtig, "Decoupled: Low-Effort Noise-Free Execution on Commodity Systems", ROSS 2016, June 2016, Kyoto, Japan







DECOUPLING: BSP



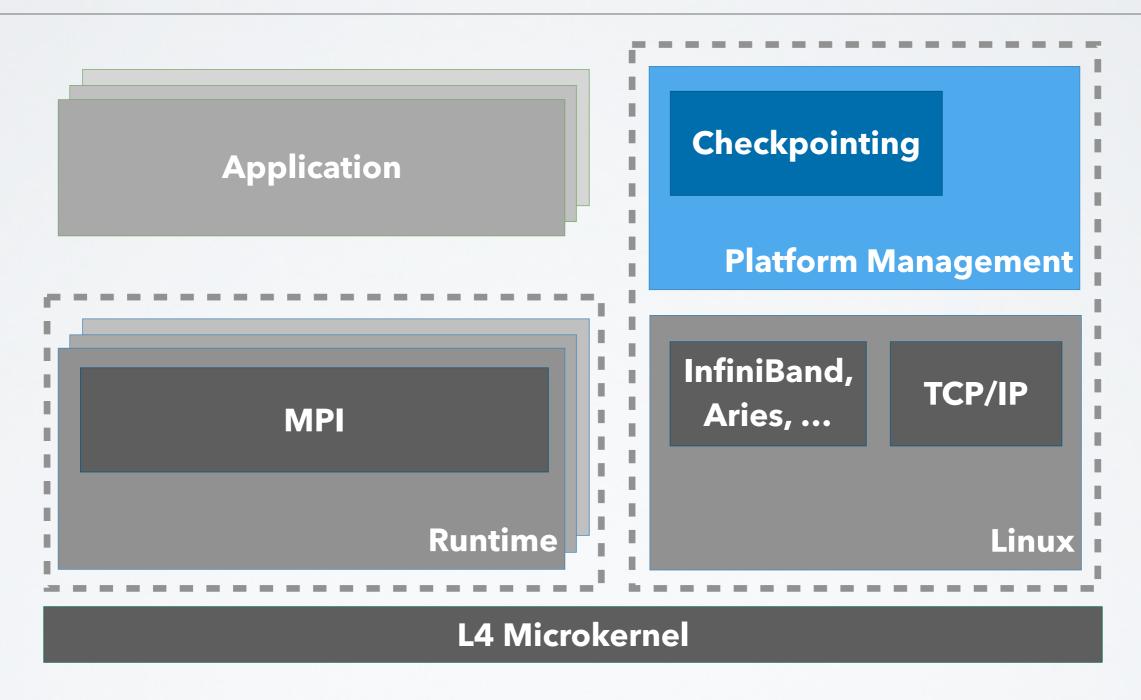
Adam Lackorzynski, Carsten Weinhold, Hermann Härtig, "Decoupled: Low-Effort Noise-Free Execution on Commodity Systems", ROSS 2016, June 2016, Kyoto, Japan







NODE ARCHITECTURE

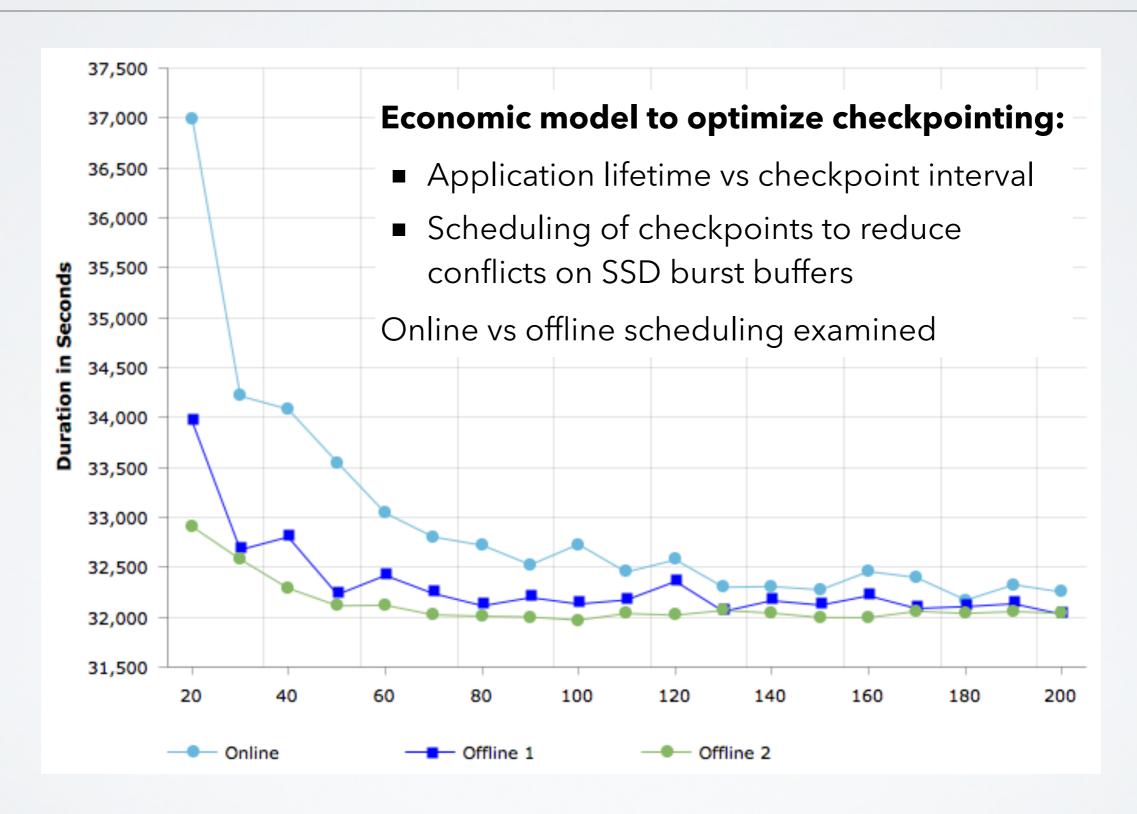








CHECKPOINTING

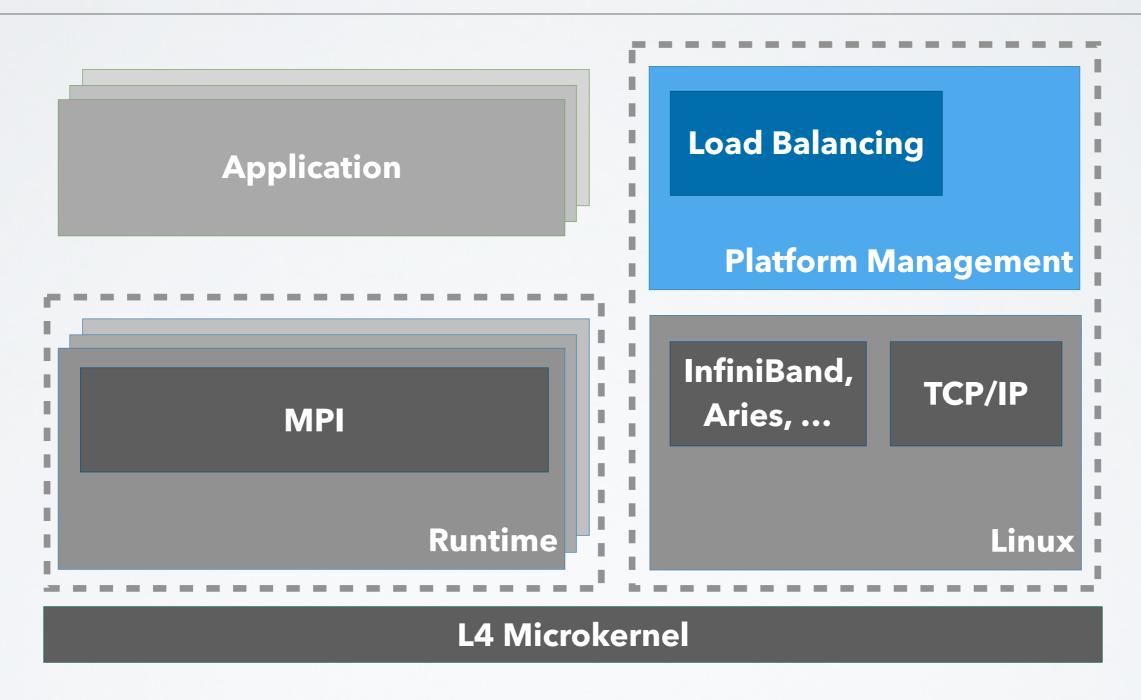








NODE ARCHITECTURE

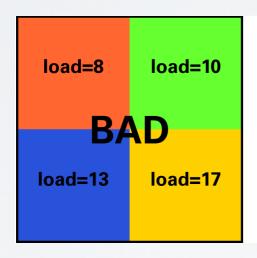


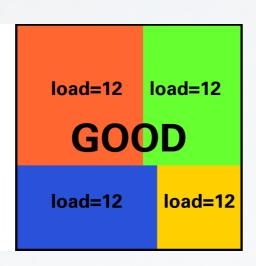






Balance workload



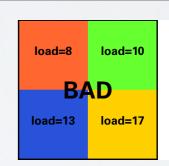


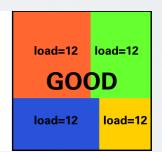




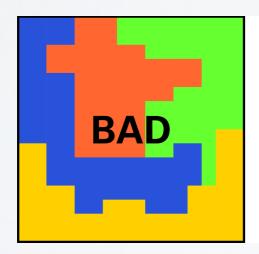


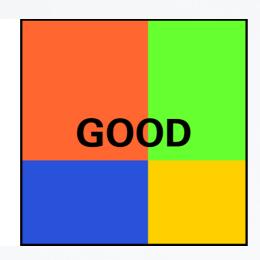
Balance workload





 Minimize communication between partitions







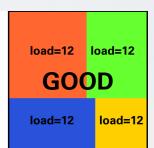




Balance workload

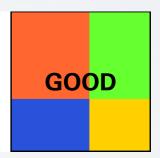
Ioad=8 Ioad=10

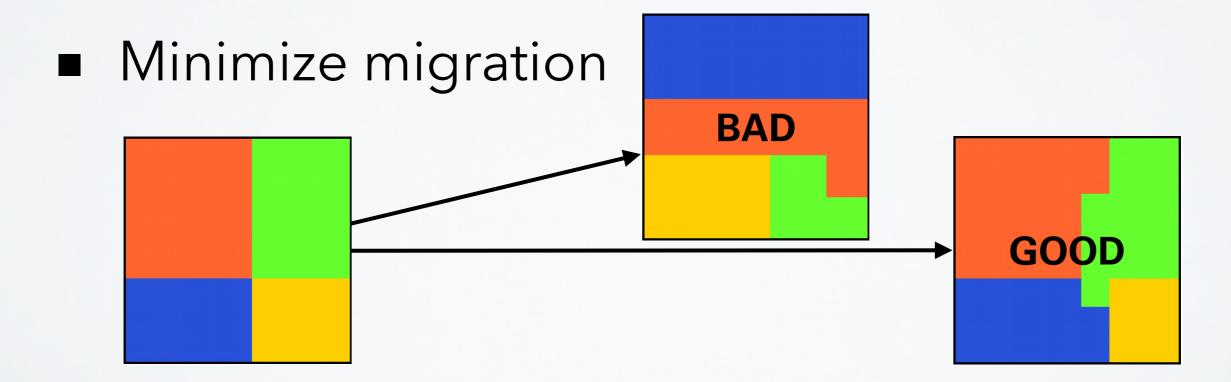
BAD
Ioad=13 Ioad=17



 Minimize communication between partitions









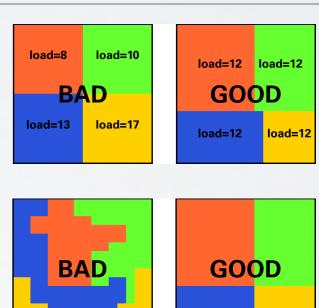


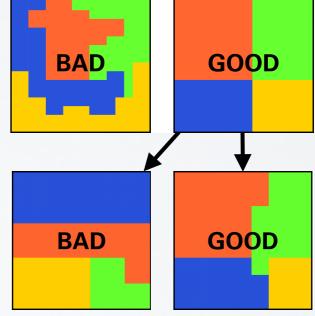


Balance workload

 Minimize communication between partitions

Minimize migration





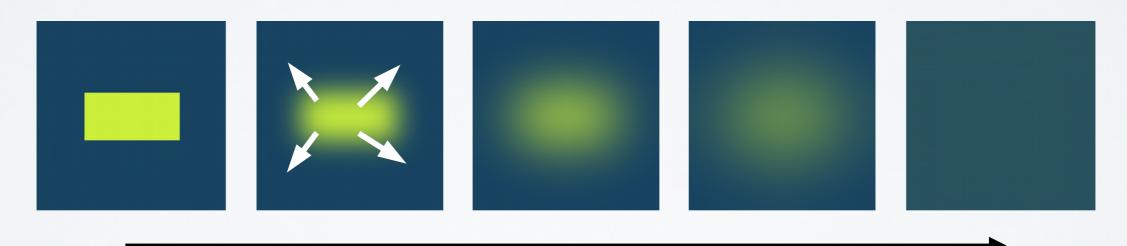
Compute new partitions fast







DIFFUSION



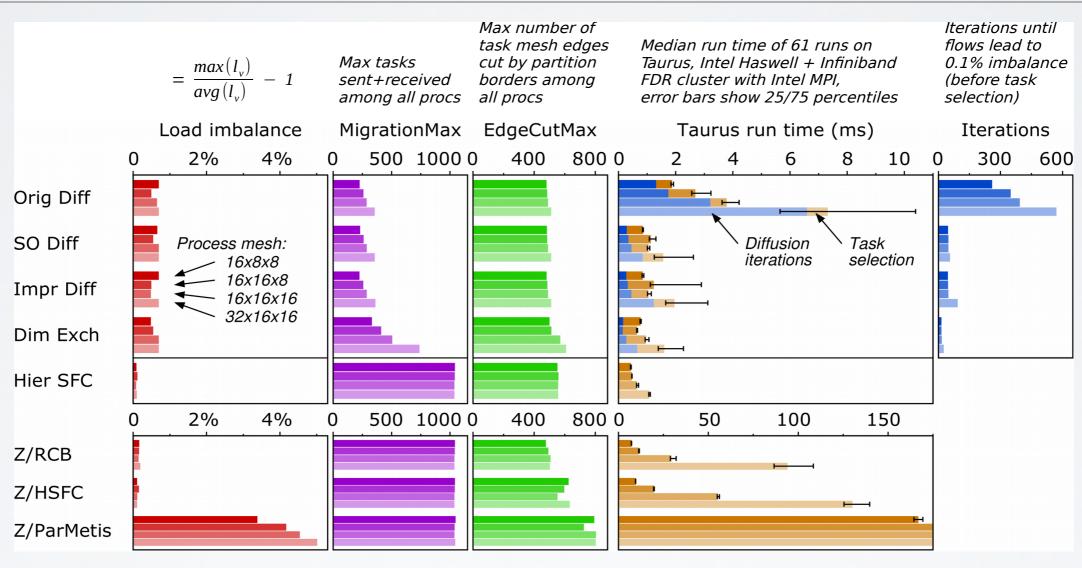
Load per node over iterations







DIFFUSION RESULTS



Diffusion leads to smallest migration

Diffusion achieves very good edge cut

Diffusion run time ~2 ms for 8192 processes, Zoltan much slower

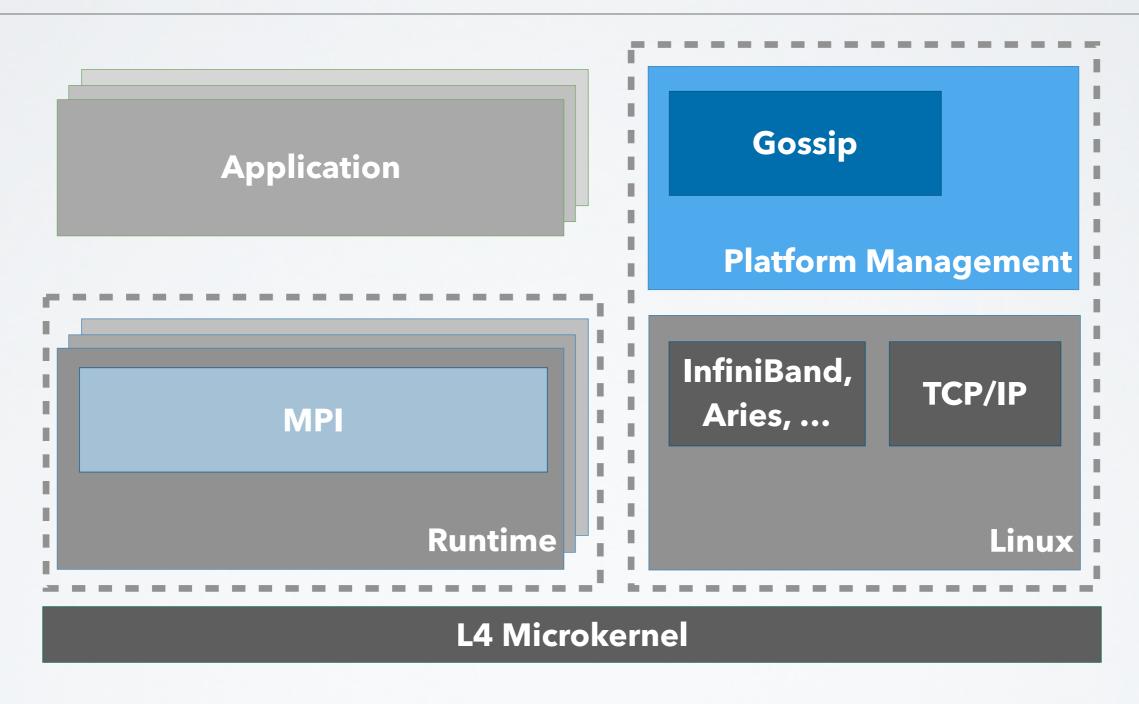
Matthias Lieber, Kerstin Gößner, Wolfgang E. Nagel, "The Potential of Diffusive Load Balancing at Large Scale", EuroMPI 2016, June 2016, Edinburgh, United Kingdom







NODE ARCHITECTURE

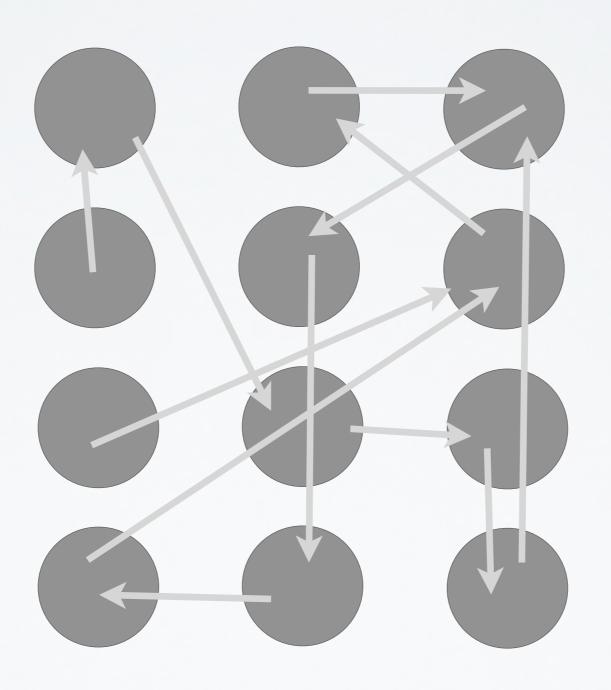








GOSSIP

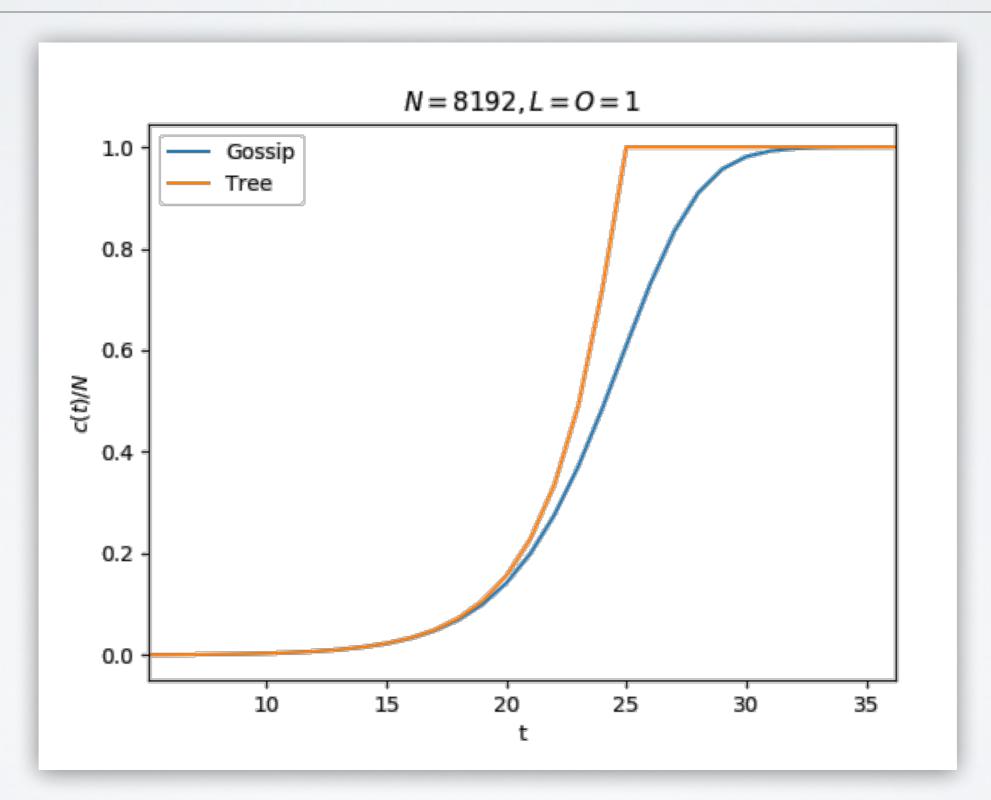








GOSSIP

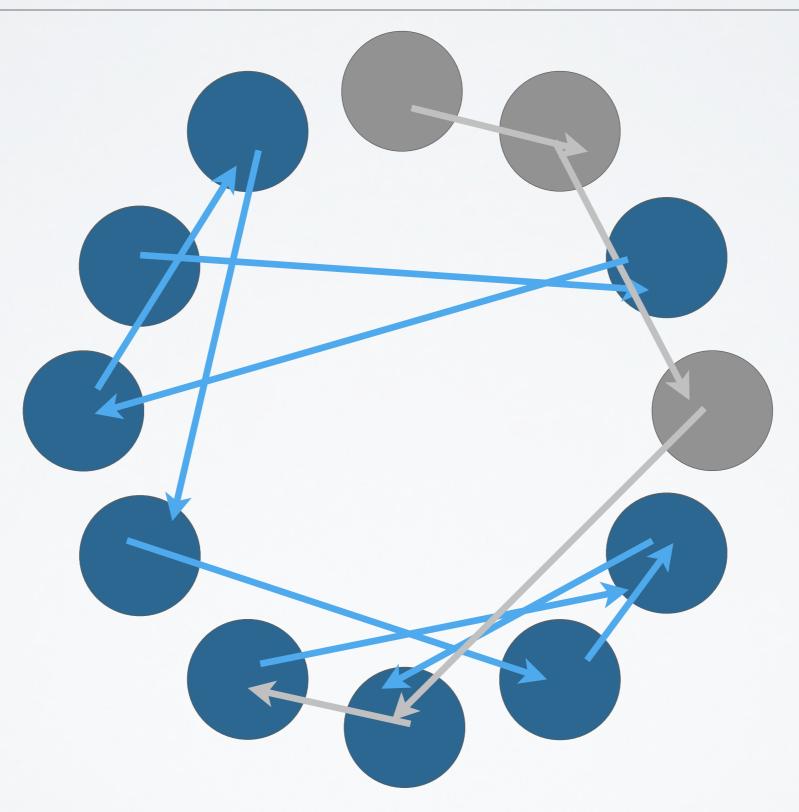








STEP 1: GOSSIP



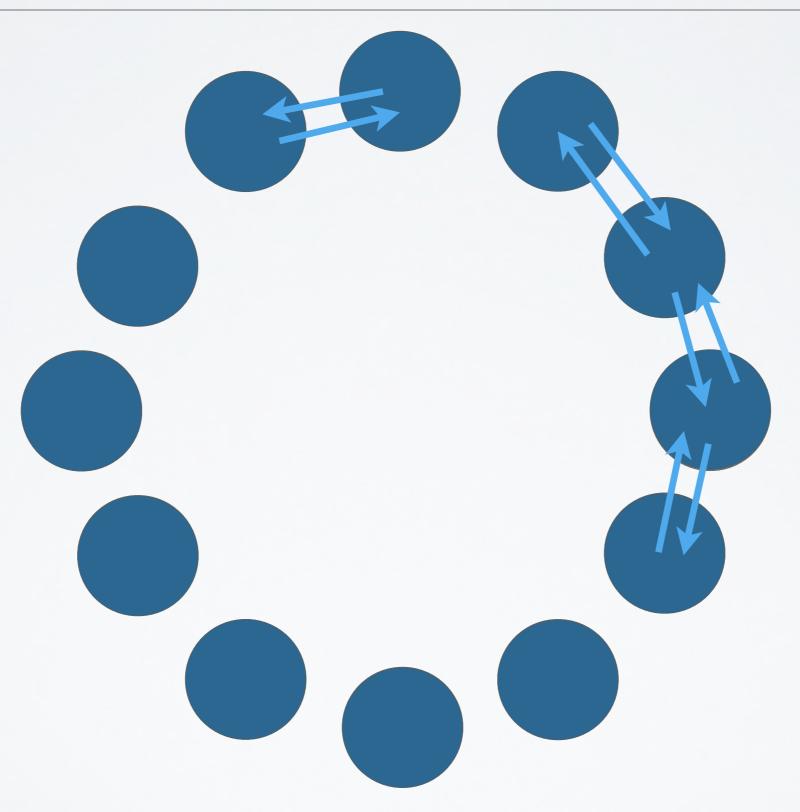
FFMK: Building an Exascale Operating System







STEP 2: CORRECTION









FT COLLECTIVES

- Two-stage algorithm: gossip + correction
- Main advantage: scalability and resilience (continues to work in presence of failures)
- Works for: fault-tolerant broadcast
- Next step: extend to operations that include barrier semantics
- Future: use in MPI?







SUMMARY

- Decoupled threads: reduced noise
- Checkpointing: Economic model
- Diffusion: may be efficient alternative
- Corrected Gossip: fault-tolerant broadcast
- Work in progress: integrate monitoring +
 gossip + decision making + migration





German Priority Programme 1648

Software for Exascale Computing